## Quantum Computing: Transforming the Digital

Age Krysta Svore Quantum Architectures and Computation (QuArC) Microsoft Research

**CRA Snowbird 2014** 



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SANTA BARBARA, Calif. - Modern computers are not unlike the looms of the industrial revolution: They follow programmed instructions to weave intricate patterns, With a loom, you see the result in a cloth or carpet. With a computer, you see it on an electronic display.

Now a group of physicists and computer scientists funded by Microsoft is trying to take the analogy of intervolven threads to what some believe will be the next great leap in competing. so-called quantum computing.

If the scientists are right, their research could lead to the design of computers that are far more powerful than today's supercomputers and could solve problems in fields as diverse as chemistry, material science, artificial intelligence and codebreaking.



Michael Freedman, Jiankas Das Sarma, and Chetan Nayals propriet a computing model in-2018 that can be used to invariant gabin, the burnlation of guarante compating. Easty hart by The Inni-Yark States'





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## Computers have come a long way



Antikythera mechanism (100 BC)



Babbage's Difference Engine (proposed 1822)



**ENIAC** (1946)



**Sequoia** (2012)



**Quantum** (2025?)

## Success of Digital Computers and Moore's Law



## Success of Digital Computers and Moore's Law



# Is there anything we can't solve on digital computers?



## Some problems are hard to solve



#### **Ultimate goal:**

Develop quantum algorithms whose complexity lies in BQP\P

## Quantum Magic: Interference



interference
 pattern =
 quantum
 coherence

Classical objects go *either* one way or the other. Quantum objects (electrons, photons) go *both* ways. Gives a quantum computation an inherent type of parallelism!

## Quantum Magic: Qubits and Superposition

 $|\psi\rangle = |0\rangle + |1\rangle = |$ 



Information encoded in the state of a two-level quantum system

 $|\psi\rangle =$ 

0 1 0 0 0 0 0 0 0 0 0 0 1 ы п п н е 1

Thanks to Charlie Marcus



## Quantum Magic: Entanglement

E







## Quantum Magic: Entanglement



 $|\psi\downarrow total\rangle = (\alpha\downarrow 0 \ |0\rangle + \beta\downarrow 0 \ |1\rangle) \otimes \ (\alpha\downarrow 1 \ |0\rangle + \beta\downarrow 1 \ |1\rangle) \otimes ... \otimes \ (\alpha\downarrow N-1 \ |0\rangle + \beta\downarrow N-1 \ |1\rangle)$ 

State of *N* non-interacting qubits: ~ *N* bits of info

Thanks to Rob Schoelkopf

## Quantum Magic: Entanglement

General state of *winteracting* qubits

32 distinct amplitudes!

1/1/2

Simulating a 200-qubit interacting system requires  $\sim_{10760}$  classical bits!

 $|\psi \downarrow total\rangle = c \downarrow 0 / 00 ... 0 + c \downarrow 1 / 00 ... 1 + ... c \downarrow 2 \uparrow N - 1 / 11 ... 1 \rangle$ 

State of *N* interacting qubits: ~ 2*tN* bits of info!

Thanks to Rob Schoelkopf

## Quantum Magic: What's the catch?





Thanks to Rob Schoelkopf

## Quantum Gates: Digital quantum computation

Basic unit: **bit** = 0 or 1 Computing: **logical** operation Basic unit: **qubit** = unit vector  $\alpha/0 + \beta/1$ 

Computing: unitary operation



**ΝΟΤ** [**-**0&1@1&0 ][**-**α@β]=[**-**β@α]



## Quantum Gates: Digital quantum computation

Basic unit: **bit** = 0 or 1 Computing: **logical** operation Description: **truth table** 



Basic unit: **qubit** = unit vector  $\alpha(0)+\beta(1)$ Computing: **unitary** operation Description: **unitary matrix** 



**[■**1&0&0&0@0&1&0&0@0&0&0

**CNOT** gate

## Quantum power unleashed: super-fast FFT



Quantum FFT



 $\# ops = \log N$ 

Example: 1GB of data = 27 ops (!!!)

## Any other catches?

#### No-cloning principle



#### I/O limitations



Quantum information cannot be copied

Input: preparing initial state can be costly Output: reading out a state is probabilistic Requirements for Quantum Computation Quantum Algorithms: Design real-world quantum algorithms for small-, medium- and large-scale quantum computers

Quantum hardware architecture: Architect a scalable, fault-tolerant, and fully programmable quantum computer

#### **Quantum software architecture:**

Program and compile complex algorithms into optimized, target-dependent (quantum and classical) instructions

## Quantum Algorithm "Wins"

Shor's Algorithm (1994)	<ul> <li>Breaks RSA, elliptic curve signatures, DSA, El-Gamal</li> <li>Exponential speedups</li> </ul>	
Solving Linear Systems of Equations (2010)	<ul> <li>Applications shown for electromagnetic wave scattering</li> <li>Exponential speedups</li> </ul>	
Quantum simulation (1982)	<ul> <li>Simulate physical systems in a quantum mechanical device</li> <li>Exponential speedups</li> </ul>	

# Cryptography







### $1387 = 19 \times 73$

=  $\times$   $\blacksquare$ 

1807082088687 4048059516561 6440590556627 8102516769401

3968599455344994595978646735

Example: (n=2048 bits) classically ~7x10<sup>15</sup> years quantum ~100 seconds

5551572 999044543 99

## signatures

## **Classical:**

 $O(\exp(n^{1}/3 (\log n)^{2}/3))$ 

## Quantum:

 $O(n \uparrow 2 \log n)$ 





VIS



## How does quantum factoring work? [Shor'94]



## Machine learning

### Solving linear systems of equations

#### Ax=b

#### is , then it requires time O(N13

## Matrix inversion can be expensive

Large inversion problems are expensive.



Seismic Processing

If you have a billion pieces of data (and sparse A) roughly 10718 operations are required.



They are ubiquitous in science and engineering.



Tomography

Quantum Matrix Inversion Algorithms: roughly 10,000 operations.



## How does the algorithm work

Steps for finding  $\mathcal{X}$  [Harrow, Hassidim, Lloyd PRL'10]  $A = \sum j \uparrow a \downarrow j \not v \downarrow j \uparrow v \downarrow j \uparrow *$ , where  $v \downarrow j$  are eigenvectors. Evolve  $e \uparrow -iAt \mid b \rangle$ .



## You can't always get what you want

Quantum computing doesn't give *x* efficiently, but allows you to sample from *x*.



Quantum least square fitting side steps this problem (Wiebe, Braun, Lloyd 2012).



Electromagnetic scattering problems can be solved this way. (Clader et al, 2013)



You have to know the right questions to ask a quantum computer.



## **Quantum Machine Learning Algorithms**

Translating classical algorithms is usually not the best approach:

- 1. You have to load all the data (at least linear time)
- 2. You have to process the data (may be exponentially faster)
- 3. You get to read-out one number as an answer (which is probabilistic)
- 4. Want another answer? Go back to step 1
# Quantum simulation

### What does quantum simulation do?

### **Physical Systems**

Quantum Chemistry



#### Superconductor Physics



#### Quantum Field Theory



#### **Computational Applications**

Emulating Quantum Computers



#### Linear Algebra



#### **Differential Equations**



### Quantum simulation

Particles can either be spinning clockwise (down) or counterclockwise (up)



There are 275 possible orientations in the quantum distribution. Cannot store this in memory for 100 particles.

### How does simulation work? [Lloyd Science'96]



### Quantum Simulation for Quantum Chemistry Ultimate problem:

Simulate molecular dynamics of *larger* systems or to *higher accuracy* 

Want to solve system *exactly* 

#### **Current solution:**

33% supercomputer usage dedicated to chemistry and materials modelingRequires simulation of exponential-size Hilbert spaceLimited to 50-70 spin-orbitals classically

#### **Quantum solution:**

Simulate molecular dynamics using *quantum simulation* Scales to 100s spin-orbitals using only 100s qubits Runtime recently reduced from O(NT11) to O(NT4) - O(NT11)



# Quantum Chemistry

#### $H = \sum pq^{\text{min}} h \downarrow pq \ a \downarrow p^{\text{min}} a \downarrow q + 1/2 \sum pqrs^{\text{min}} h \downarrow pqrs \ a \downarrow p^{\text{min}}$

Can quantum cherm Chemistry: M. B.computer: Dave We The Trotter Step Size Required for Accurate Quantum Simulation of Quantum Chemistrym Chemistry: M. B.Hastings, Matthias T David Poulin, M. B. Hastings, Dave Wecker, Nathan Wiebe, Andrew C. Doherty, Matthiasm Chemistry: M. B.

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Ferredoxin (*Fel2 Sl2*) used in many metabolic reactions including energy transport in photosynthesis

- > Intractable on a classical computer
  - *First paper: ~300 million years to solve*
- Second paper: ~30 years to solve (10 77 reduction)

Third paper: ~300 seconds to solve (another 1073 reduction)

quantum chemistry protetyvechilisa, Hvverili4a} which relied on exponentially costly classical exact simulation. be needed.

http://arxiv.org/abs/1406.4920

otter-Suzuki based on a quantum nations are nic in the not require ny operations in the parallel increase in rder in the error at given e Hamiltonian to timestep. All of on and detailed

## Quantum Chemistry

#### $H = \sum pq^{\text{min}} h \downarrow pq a \downarrow p^{\text{min}} a \downarrow q + 1/2 \sum pqrs^{\text{min}} h \downarrow pqrs a \downarrow p^{\text{min}}$



# **Application: Nitrogen Fixation**

#### **Ultimate problem:**

Find catalyst to convert nitrogen to ammonia at room temperature

Reduce energy for conversion of air to fertilizer

#### **Current solution:**

Uses Haber process developed in 1909 Requires high pressures and temperatures Cost: 3-5% of the worlds natural gas production (1-2% of the world's annual energy)

#### **Quantum solution:**

~ 100-200 qubits: Design the catalyst to enable inexpensive fertilizer production



# **Application: Carbon Capture**

#### **Ultimate problem:**

Find catalyst to extract carbon dioxide from atmosphere Reduce 80-90% of emitted carbon dioxide

#### **Current solution:**

Capture at point sources Results in 21-90% increase in energy cost

#### **Quantum solution:**

~ 100-200 qubits: Design a catalyst to enable carbon dioxide extraction from air



### How much faster is it?

Imagine you have 100 interacting electrons in a superconductor



Quantum computing makes testing models of high-temperature superconductivity conceivable.

Supercomputer simulation requires roughly 10*1*34 operations.



Quantum simulation requires roughly 10710 operations. [Wiebe,Berry, Hoyer, Sander JPA'11]

Wiebe,Childs, QIC'12]



Exact gate counts can be found using [Raeisi, Wiebe, Sanders NJP'12], [Wecker, Bauer, Clark, Hastings, Troyer arXiv'14].

## Quantum Algorithm Opportunities

Quantum simulation	<ul> <li>Extend q. chem. method to solid state materials</li> <li>E.g., high temp. superconductivity</li> <li>~ 2000 qubits; linear or quad. scaling</li> </ul>	
Machine learning	<ul> <li>Clustering, regression, classification</li> <li>Polynomial speedups to date</li> <li>Can we harness interference to produce better inference models?</li> </ul>	0.9 0.8 0.7 0.6 0.5 0.4 0.3 0.2 0.4 0.6 0.8
Cryptography	<ul> <li>RSA, DSA, elliptic curve signatures, El-Gamal</li> <li>What questions should we pose to a quantum computer?</li> </ul>	HICCL SE CON

### **Quantum Hardware Technologies**

#### lon traps





### NV centers



Quantum dots



#### Linear optics



### Topological





### Ages of Quantum Computing

"Age of Qu. Error Correction" "Age of Qu. Feedback" "Age of Measurement" "Age of Entanglement" "Age of Coherence"



M. Devoret and R. Schoelkopf, Science (2013)

Thanks to Rob Schoelkopf

### **Classical Error Correction**



Probability *p* of having a bit flipped

#### **Repetition code**: redundantly encode, majority voting $0 \rightarrow 000$ $1 \rightarrow 111$

Reduces classical error rate to  $3p^2 - 2p^3$ 

Can we do this for quantum computing? Some reasons to think no:

- "No cloning" theorem
- Errors are continuous (or are they?)
- Measurements change the state

Thanks to Rob Schoelkopf

### Different Error Correction Architectures Standard QEC Surface Code



## Overhead required in known schemes: 1,000 – 10,000 actual qubits for every logical!!

(+ concatenation!)

- threshold ~  $10^{-4}$
- many ops., syndromes per QEC cycle

- threshold ~ 1%
- large system to see effects?
- good local gates (10<sup>-4</sup>?) remote gates fair (90%?)
- then construct QEC as software layer?

Thanks to Rob Schoelkopf



### The man and his particle

Chi l'ha visto?



sedimaria da fisica. toories all' Univeraria di Napoli, e misteriosamente scomparso dagli ultinui di marzo. Di anni 31, alto metri 1.70, snello, con capelli neri, orchi acuri, una lunga encatrice sul dorso di una mano, Chi nesapesse qualcosn è pregato di scrivere al R. P. E. Maria-

necci. Viale Regina Margherita 46 Roma

1938

Image courtesy of Leo Kouwenhoven



→ Electric charge is zero
 → Energy is zero
 → Everything is zero (except mass)

How to measure the "niks" particle?
→ That's why it was not yet detected!

Image courtesy of Leo Kouwenhoven

#### New directions in the pursuit of Majorana fermions in solid state systems

Jason Alicea<sup>1</sup>

<sup>1</sup>Department of Physics and Astronomy, University of California, Irvine, California 92697 (Dated: February 8, 2012)



### **Topology provides natural immunity to noise!**



FIG. 6. (a) Basic architecture required to stabilize a topological superconducting state in a 1D spin-orbit-coupled wire. (b) Band structure for the wire when time-reversal symmetry is present (red and blue curves) and broken by a magnetic field (black curves). When the chemical potential lies within the field-induced gap at k = 0, the wire appears 'spinless'. Incorporating the pairing induced by the proximate superconductor leads to the phase diagram in (c). The endpoints of topological (green) segments of the wire host localized, zero-energy Majorana modes as shown in (d).



# Non-Abelian statistics and topological quantum information processing in 1D wire networks

Jason Alicea<sup>1\*</sup>, Yuval Oreg<sup>2</sup>, Gil Refael<sup>3</sup>, Felix von Oppen<sup>4</sup> and Matthew P. A. Fisher<sup>3,5</sup>



# Epitaxial growth of InAs (or InSb) nanowires



#### Signatures of Majorana Fermions in Hybrid Superconductor-Semiconductor Nanowire Devices

V. Mourik,<sup>1\*</sup> K. Zuo,<sup>1\*</sup> S. M. Frolov,<sup>1</sup> S. R. Plissard,<sup>2</sup> E. P. A. M. Bakkers,<sup>1,2</sup> L. P. Kouwenhoven<sup>1</sup><sup>†</sup>



### Hardware provides error correction: Only ~10-100s for every logical???



## LIQ*Ui*): Quantum Software Architecture

- Enables easy programming and simulation of complex quantum circuits
- Allows retargeting of circuits for various purposes: simulation, rendering, optimization, noise modeling, and export
- Provides state-of-the-art quantum circuit simulation tools



The LIQ*Ui*|> *platform* [Wecker, Svore, 2014]

### Quantum Gates Hamiltonian $H=\ln(U)/\Delta t$

Туре	Basis	U	Name	Sym	Туре	Basis	U	Name	Sym
Pauli	{/0 <i>)</i> ,  1}}	[ <b>■</b> 0&1@1& 0]	Х	-X		{ <b>—</b>	[ <b>■</b> 1&0&0 &0 <i>@</i> 0&1		-
	{/0 <i>)</i> ,  1}}	[ <b>■</b> 0& − <i>i@i</i> &0]	Y	-Y	Controlled Not	00),/ 01 <i>)</i> , / 10 <i>)</i> ,	&0&0@0 &0&0&1	CNOT (CX)	
Z Rotation	{/0 <i>}</i> ,  1}}	[ <b>■</b> 1&0@0& -1]	Z			11} }	@0&0&1 &0]		*
<i>e</i> îπ/2	{/0 <i>}</i> ,  1\}	[ <b>■</b> 1&0@0&	S	$\begin{bmatrix} 3 \\ -T \end{bmatrix}$		{ <b>■</b> / 00}/	[ <b>■</b> 1&0&0 &0 <i>@</i> 0&0		- <del>*</del> -
	1) <i>}</i> {/0 <i>)</i> ,  1) <i>}</i>	τ] [ <b>■</b> 1&0@0& eîiπ/4]	т	- <u>R4</u> -		01), / 10),   11} }	$\&1\&0@0\\\&1\&0\&0\\@0\&0\&0\\$	SWAP	
	{/0 <i>)</i> ,  1}}	[ <b>■</b> 0&1@1& eîiπ/8 ]	R4		Moasuro	{/0 <i>}</i> ,	&1 J	N/	
Identity	<i>{/</i> 0 <i>},</i>	[■1&0@0&	I	-H	weasure	$1\rangle\}$	QUDIT TO BIT	IVI	$=  0\rangle$
identity	1)}	$\frac{1}{1/\sqrt{2}}$			Binary Control	{/0 <i>)</i> ,  1>}	Conditional Application	BC	

## Shor's algorithm: 4 bits $\cong$ 8200 gates

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**Circuit for Shor's algorithm using 2n+3 qubits** – Stéphane Beauregard



## Simulating Shor's Algorithm



# L[QUi] for Compilation onto Hardware

```
let QFT (qs : Qs)
    let n = qs.Length - 1
    for i = 0 to n do
    let q = qs.[i]
    H q
    for j = (i + 1) to n do
        let theta = 2.0 * Math.Pl /
        float(1 <<< (j - i + 1))
        CRz theta qs.[j] q
    for i = 0 to ((n - 1) / 2) do
    SWAP qs.[i] qs.[n - i]
```

let QftOp = compile QFT
let QftOp' = adjoint QftOp



### **Spin-Glass Models**

#### $H(t) = \Gamma(t) \sum_{i=1}^{N} \Delta_{i} \sigma_{i} + \Lambda(t) \sum_{i=1}^{N} h_{i} \sigma_{i}$

Quantum annealir Sergio Boixo, Troel David Wecker, Dar

Quantum technolod devices, such as qu random number ge with capabilities ex annealer, in particu evolving a known ir towards the ground problem. Here, we D-Wave One devic strong correlations annealer, in contra and classical annea that the device per evidence for quant level crossings cha computational pow classical algorithms



# Yes, you can

### Sciencexpress

### Defining and detecting quantum speedup

Troels F. Rønnow,<sup>1</sup> Zhihui Wang,<sup>2,3</sup> Joshua Job,<sup>3,4</sup> Sergio Boixo,<sup>5</sup> Sergei V. Isakov,<sup>6</sup> David Wecker,<sup>7</sup> John M. Martinis,<sup>8</sup> Daniel A. Lidar,<sup>2,3,4,9</sup> Matthias Troyer<sup>1</sup>\*

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\*Corresponding author. E-mail: troyer@phys.ethz.ch

The development of small-scale quantum devices raises the question of how to fairly assess and detect quantum speedup. Here we show how to define and measure quantum speedup, and how to avoid pitfalls that might mask or fake such a speedup. We illustrate our discussion with data from tests run on a D-Wave Two device with up to 503 qubits. Using random spin glass instances as a benchmark, we find no evidence of quantum speedup when the entire data set is considered, and obtain inconclusive results when comparing subsets of instances on an instance-by-instance basis. Our results do not rule out the possibility of speedup for other classes of problems and illustrate the subtle nature of the quantum speedup question.

#### F. Rønnow, Zhihui David Wecker, John

uantum devices are the es, and of how to ne and measure avoid pitfalls that our discussion with e Two device with device on random ed classical and speedup when the e results when stance basis. Our he possibility of that quantum posed.

#### <u>401.2910</u>

# Conclusions

Quantum computers exploit interference and superposition to solve problems.

How big/fast does a quantum computer have to be to have an advantage?

[Boixo, Ronnow et al '13] [Wecker, Bauer et al '14]

Exponential speedups for *certain* simulation, cryptography, linear algebra problems.

How do you compile, test, and debug quantum algorithms?

[Wiebe, Kliuchnikov'13] [Bocharov, Gurevich, Svore'13] [Wecker, Svore Geller' 14] quantum computer? [Wiebe, Braun, Lloyd '12]

[Wiebe, Grenade et al '13]

What are the right

questions to ask a

What other problems does a quantum computer solve better or faster?



# Station Q



#### University Partners David Reilly U. Sydney Amir Charlie Leo Yacoby Kouwenhoven Marcus Harvard NBI Delft Chris Dale Mike **Palmstrom** Van Harlingen Manfra UCSB UIUC Purdue Matthias Sankar Bert Halperin Das Sarma Troyer **ETH Zurich** Harvard **U** Maryland

# Outlook

Training of students

Field is becoming a merging and now entering age to be treated as a computer science Encourage more funding

Need research at all levels and in many areas

We have qubits, we have algorithms, now we need the computer science and corresponding software infrastructure

Need more students in quantum algorithms --- appification of quantum computing!

Support of quantum computing as "quantum computer science"

Not just as physics

### http://research.microsoft.com/groups/quarc/

### http://research.microsoft.com/en-us/labs/stationq/





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